

## ECE 259A: Midterm Exam

**Instructions:** There are four problems, weighted as shown below. The exam is open book and open notes: you may use any auxiliary material that you like as long as it is on paper.

**Good luck!**

**Problem 1.** (25 points)

Let  $\mathbb{C}$  be the binary linear code of length  $n = 10$ , generated by the following matrix

$$G = \begin{bmatrix} 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 & 1 & 0 \\ 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 & 1 \\ 0 & 0 & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 0 \\ 0 & 0 & 0 & 1 & 1 & 1 & 1 & 0 & 0 & 0 \end{bmatrix}$$

- a. Express  $G$  in systematic form, and compute a parity-check matrix for this code.
- b. What is the minimum distance of  $\mathbb{C}^\perp$ , the dual code of  $\mathbb{C}$ ?
- c. Compute the syndrome of the vector  $\underline{y} = (1, 0, 1, 1, 0, 0, 1, 1, 0, 0)$  with respect to the parity-check matrix found in part (a).
- d. If the vector  $\underline{y} = (1, 0, 1, 1, 0, 0, 1, 1, 0, 0)$  from part (c) is observed at the output of a binary symmetric channel, what is the most likely transmitted codeword of  $\mathbb{C}$ ?

Hint: Compare the syndrome from part (c) with the columns of the parity-check matrix.

**Problem 2.** (20 points)

A binary code  $\mathbb{C}$  is said to be *doubly-even* if the weight of all the codewords of  $\mathbb{C}$  is divisible by 4. Suppose that a binary linear code  $\mathbb{C}$  and its dual  $\mathbb{C}^\perp$  are both doubly-even. Prove that  $\mathbb{C} = \mathbb{C}^\perp$ .

**Problem 3.** (25 points)

Let  $\mathcal{E}$  be an arbitrary subset of  $\mathbb{F}_2^n$  of cardinality  $|\mathcal{E}| = M$ , such that  $\underline{0} \in \mathcal{E}$ . We say that a code  $\mathbb{C} \subset \mathbb{F}_2^n$  detects all error patterns in the set  $\mathcal{E}$  if  $\underline{c}_1 + \underline{e} \neq \underline{c}_2$  for all  $\underline{e} \in \mathcal{E}$  and for all distinct  $\underline{c}_1, \underline{c}_2 \in \mathbb{C}$  (or, equivalently, if  $(\underline{c} + \mathcal{E}) \cap \mathbb{C} = \{\underline{c}\}$  for all  $\underline{c} \in \mathbb{C}$ ).

Show that there exist codes of cardinality  $|\mathbb{C}| \geq 2^n / M$  that detect all error patterns in the set  $\mathcal{E}$ .

**Problem 4.** (30 points)

Let  $\mathbb{C}$  be an  $(n, k, d)$  linear code over the field  $\text{GF}(q) = \{0, \alpha^0, \alpha^1, \dots, \alpha^{q-2}\}$ , where  $q \geq 3$ . It is known that  $\mathbb{C}$  is at the same time an MDS code and a non-trivial perfect code.

- a. What are the parameters  $n, k, d$  of  $\mathbb{C}$  (in terms of  $q$ )?
- b. Write down a parity-check matrix for  $\mathbb{C}$ .

Hint: Use the fact that the only non-trivial perfect linear codes over a field are the binary Golay code over  $\text{GF}(2)$ , the ternary Golay code over  $\text{GF}(3)$ , and the Hamming codes over  $\text{GF}(q)$  for all  $q$ .